## The Nash bargaining solution:

This is referred to the axiomatic approach to bargaining. The reason is that Nash was interested in finding a solution with particular properties to a bargaining problem. Let us first define the "problem" and the properties that the "solution" must have. There are two agents<sup>1</sup>, j = 1, 2, with utility functions  $u_j$ . There is an arbitrary set of outcomes A. D is the outcome in case the agents cannot reach an agreement (disagreement or threat point). Define  $S = \{(u_1(a), u_2(a)), a \in A\}$  and  $d = (d_1, d_2)$ , where  $d_j = u_j(D)$ . Suppose that S is compact and convex and that  $d \in S$ . Also assume that  $\exists s \in S$ , such that  $s_j > d_j$ ,  $\forall j = 1, 2$ .

(S, d) is the bargaining problem

 $f : (S,d) \longrightarrow S$  is a solution to (S,d)

Nash was looking for a "solution" with the following properties:

• (A1) Invariance to utility choices:

Given (S, d) and (S', d') defined by  $s'_j = \alpha_j s_j + \beta_j$  and  $d'_j = \alpha_j d_j + \beta_j$ , then  $f_j(S', d') = \alpha_j f_j(S, d) + \beta_j$ 

• (A2) Symmetry:

If  $d_1 = d_2$  and  $(s_1, s_2) \in S \iff (s_2, s_1) \in S$ , then  $f_1(S, d) = f_2(S, d)$ 

- (A3) Independence of irrelevant alternatives:
- If (S, d) and (S', d) satisfy  $S \subset S'$  and  $f(S', d) \in S$ , then f(S, d) = f(S', d)
- (A4) Pareto efficiency:

Given (S, d), if  $s \in S$  and  $s' \in S$  and  $s'_i > s_j, \forall j = 1, 2$ , then  $f(S, d) \neq s$ 

Nash (1950) showed that the **unique** solution to this problem  $is^2$ :

$$f(S,d) = \underset{s_1 \ge d_1, s_2 \ge d_2}{\operatorname{Arg\,max}} (s_1 - d_1) (s_2 - d_2) \tag{1}$$

To sketch the proof, it will be useful to draw a graph. By (A1), one can choose the set of possible outcome  $S_1$ , such that d = (0,0) (normalization of the utility functions). Denote by  $S_2$  the intersection of  $S_1$  and the positive quadrant. Let  $(u_1^*, u_2^*) = Arg \max_{s \in S_2} u_1 u_2$ . By assumption,  $S_2$  is non-empty, compact and convex, which guarantees existence of the maximizers. Uniqueness is obtained from the convexity assumption. By (A1), choose  $u_1, u_2$  such that  $(u_1^*, u_2^*) = (u^*, u^*)$  lies on the 45° line (normalization of the utility functions).

 $<sup>^{1}</sup>$ In what follows, we are only interested in situations where **two** agents bargain.

<sup>&</sup>lt;sup>2</sup>Remark that if requirement (A2) were dropped, then there is a continuum of solutions:

 $f_{\theta}(S,d) = \underset{s_1 \ge d_1, s_2 \ge d_2}{Arg \max} (s_1 - d_1)^{\theta} (s_2 - d_2)^{(1-\theta)}$ 

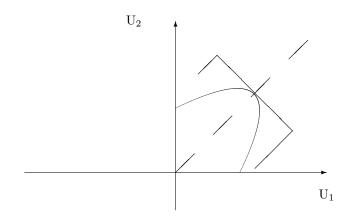


Figure 1: Graphical transformation used in Nash's proof

Notice that every point of  $S_2$  is such that  $u_1 + u_2 \leq 2u^{*3}$ . Let B be a square, symmetric relative to the 45° line, one side of which is supported by  $u_1 + u_2 = 2u^*$ , that includes S (of course, it is not unique). It exists since S is bounded. Then by (A2), f(B,O) is located on the 45° line. By (A4),  $f(B,O) = (u^*, u^*)$ . By (A3), f(S,O) = f(B,O). Hence, given the normalizations performed, f(S,d) is located at  $(u^*, u^*)$ . Remarkably, it can be proved that uniqueness of the bargaining solution cannot be obtained with a proper subset of these four axioms.

<sup>&</sup>lt;sup>3</sup>Suppose that there exists a point  $M = (u_1, u_2) \in S_2$  such that  $u_1 + u_2 > 2u^*$ . Then, there exists a point between M and  $N = (u^*, u^*)$  that belongs to S, for which  $u_1u_2 > u^{*2}$  (by convexity of  $S_2$ ).